

Classical simulation of quantum circuits with partial and graphical stabiliser decompositions

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Quantum circuit simulation

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 - ▶ direct state simulation
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 - ▶ contraction order finding
 - ▶ edge cutting
 - ▶ ...

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Broadly, two ways to do this:

- ▶ Tensor-network based methods:
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 - ▶ ...
- ▶ Stabiliser-decomposition based methods:
 - ▶ Stabiliser extent
 - ▶ Stabiliser rank ← **This talk**

Simulating using stabiliser rank

- ▶ Start with Clifford+T circuit.
- ▶ Write each T gate as magic state injection.
- ▶ Decompose T states into sum of stabilisers.
- ▶ Efficiently simulate resulting Clifford circuits.
- ▶ Add results together.
- ▶ We're done!

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What's the catch?

Stabiliser rank of k T states scales exponentially with k .

... But it's not just 2^k terms. We can do better!

Stabiliser ranks of T magic states

Recall $|T\rangle \propto |0\rangle + e^{i\pi/4}|1\rangle$.

So $\chi(|T\rangle) = 2$ and hence $\chi(|T\rangle^{\otimes k}) \leq 2^k$.

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$$\chi(|T\rangle^{\otimes k}) = \chi((|T\rangle^{\otimes 2})^{\otimes k/2}) \leq 2^{k/2} = 2^{0.5k}.$$

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Turns out $\chi(|T\rangle^{\otimes 6}) \leq 7$ so

$\chi(|T\rangle^{\otimes k}) \leq 2^{\alpha k}$ where $\alpha = \log_2(7)/6 \approx 0.467$.

Found by Bravyi, Smith, Smolin (BSS) in 2016.

Even have $\chi(|T\rangle^{\otimes 6}) = 6$ (previous talk): $\alpha \approx 0.431$.

Our idea

Do stabiliser rank decompositions,
but with *ZX-diagrams* instead of circuits!

Benefit 1: optimise intermediate ZX-diagrams to reduce T-count.

Benefit 2: Can use *fancier* stabiliser decompositions.

Spiders

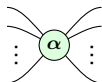
What gates are to circuits, *spiders* are to ZX-diagrams.

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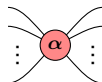
Z-spider

$$|0 \cdots 0\rangle\langle 0 \cdots 0| \\ + e^{i\alpha} |1 \cdots 1\rangle\langle 1 \cdots 1|$$





X-spider

$$|+\cdots+\rangle\langle+\cdots+| \\ + e^{i\alpha} |-\cdots-\rangle\langle-\cdots-|$$



Spiders

What gates are to circuits, *spiders* are to ZX-diagrams.

Z-spider	X-spider
$ 0 \dots 0\rangle\langle 0 \dots 0 $	$ +\dots+\rangle\langle +\dots+ $
$+e^{i\alpha} 1 \dots 1\rangle\langle 1 \dots 1 $	$+e^{i\alpha} -\dots-\rangle\langle -\dots- $
	

For example:

$$\text{---} \textcircled{\alpha} \text{---} = |0\rangle\langle 0| + e^{i\alpha}|1\rangle\langle 1| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 0 \\ 0 & e^{i\alpha} \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\alpha} \end{pmatrix}$$

$$\text{---} \textcircled{\alpha} \text{---} = |+\rangle\langle +| + e^{i\alpha}|-\rangle\langle -| = \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix} + \frac{1}{2} e^{i\alpha} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$$

Spiders cont.

If $\alpha = 0$ we drop the label:

$$\begin{array}{c} \diagup \quad \diagdown \\ \vdots \quad \vdots \\ \text{---} \circ \text{---} \\ \vdots \quad \vdots \\ \diagdown \quad \diagup \end{array} = |0 \cdots 0\rangle\langle 0 \cdots 0| + |1 \cdots 1\rangle\langle 1 \cdots 1|$$

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Example:

$$\begin{array}{l} \text{●} \text{---} = |+\rangle + |-\rangle = \sqrt{2}|0\rangle \qquad \text{○} \text{---} = |0\rangle + |1\rangle = \sqrt{2}|+\rangle \\ \text{●} \text{---} = |+\rangle - |-\rangle = \sqrt{2}|1\rangle \qquad \text{○} \text{---} = |0\rangle - |1\rangle = \sqrt{2}|-\rangle \end{array}$$

Formal composition

Spiders can be composed in two ways.

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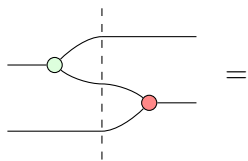
Vertical composition gives tensor product:

$$\text{---} \circlearrowleft = \begin{pmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\text{---} \circlearrowleft = \begin{pmatrix} 1 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 1 \end{pmatrix} \otimes \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

Formal composition

Horizontal composition is regular composition of linear maps:



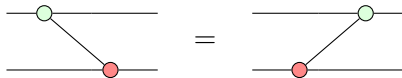
$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

Building ZX-diagrams

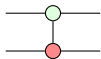
Any ZX-diagram is built by simply iterating these vertical and horizontal compositions

Symmetries

Note:

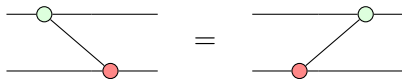


Hence, we may write

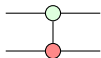


Symmetries

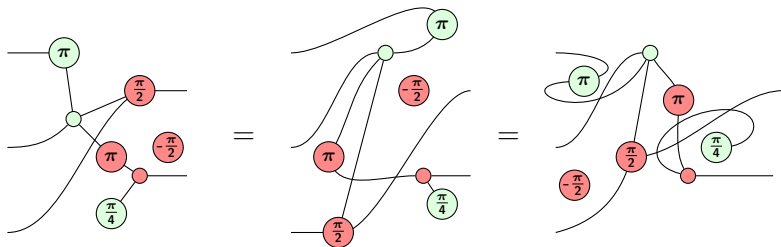
Note:



Hence, we may write



In general: *only connectivity matters*



New algorithm

- ▶ Write Clifford+T circuit as ZX-diagram.

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Let's go through these steps in more detail.

- ▶ Writing circuit as ZX-diagram
- ▶ Optimising ZX-diagram
- ▶ Decomposing magic states

Writing Clifford+T circuit as ZX-diagram

$$\text{CNOT} = \sqrt{2} \begin{array}{c} \text{---} \circ \text{---} \\ | \\ \text{---} \bullet \text{---} \end{array} \quad \text{S} = \text{---} \circ \left(\frac{\pi}{2} \right) \text{---} \quad \text{Had} = \text{---} \square \text{---}$$

$$\text{T} = \text{---} \circ \left(\frac{\pi}{4} \right) \text{---} \quad |x\rangle = \frac{1}{\sqrt{2}} \circ (x\pi) \text{---} \quad \langle x| = \frac{1}{\sqrt{2}} \text{---} \circ (x\pi)$$

$(x \in \{0, 1\})$

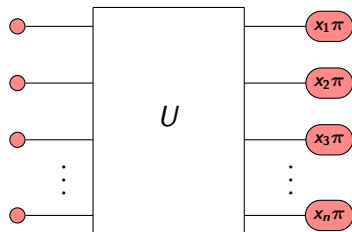
Writing Clifford+T circuit as ZX-diagram

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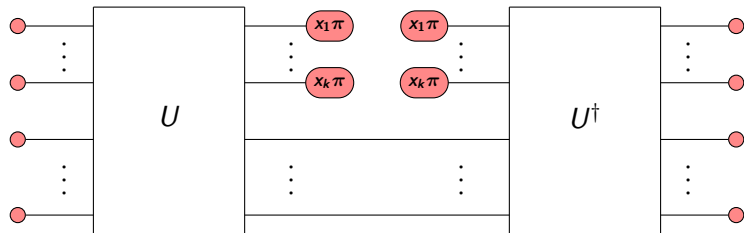
Calculating single amplitude:

$$\langle \vec{x} | U | 0 \dots 0 \rangle = \left(\frac{1}{\sqrt{2}} \right)^{2n}$$



Marginal probabilities

To calculate marginal probability, use *doubling* technique:



Strong simulation vs Weak simulation

Weak sim: approx. sample from the same output distribution.

Strong sim: approx. calculate any marginal probability.

Strong simulation vs Weak simulation

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Strong sim: approx. calculate any marginal probability.

We are doing *exact* strong simulation here.

- ▶ Writing circuit as ZX-diagram ✓
- ▶ Optimising ZX-diagram
- ▶ Decomposing magic states

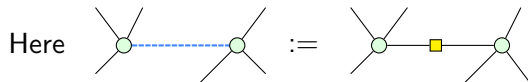
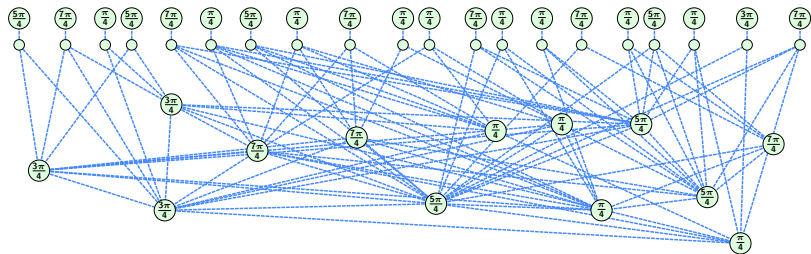
Simplifying ZX-diagrams

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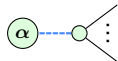
This transforms every ZX-diagram into something like this:



Properties of reduced diagram

The important part:

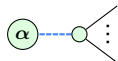
- ▶ Every spider carries a non-Clifford phase,
- ▶ or is part of a *phase gadget*:



Properties of reduced diagram

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- ▶ Every spider carries a non-Clifford phase,
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Particularly: if original circuit had k T gates, resulting diagram has $\leq 2k$ spiders (regardless of #qubits or #gates).

- ▶ Writing circuit as ZX-diagram ✓
- ▶ Optimising ZX-diagram ✓
- ▶ Decomposing magic states

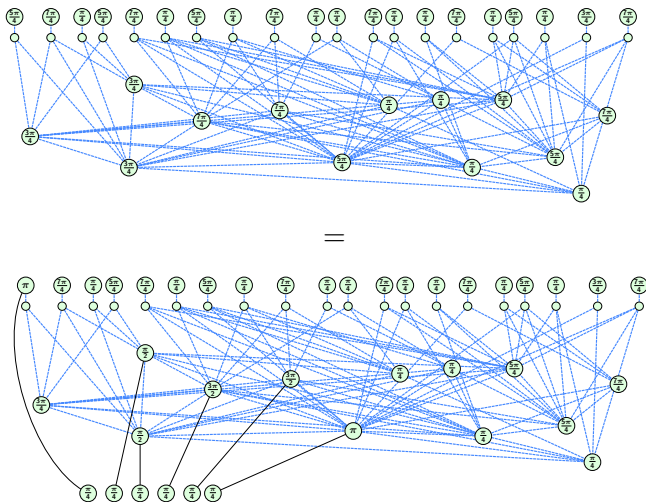
Decomposing T-like spiders

The 6-to-7 magic state decomposition in ZX is:

$$e^{i\pi/4} \begin{array}{c} | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \\ | \\ \textcircled{\frac{\pi}{4}} \end{array} = +2e^{i\pi/4} \begin{array}{c} | \\ | \\ | \\ | \\ | \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ | \\ | \\ | \\ | \end{array}$$
$$-\frac{1+\sqrt{2}}{4} \begin{array}{c} | \\ \circ \\ | \\ \circ \\ | \\ \circ \\ | \\ \circ \\ | \\ \circ \\ | \\ \circ \end{array} + \frac{1-\sqrt{2}}{4} \begin{array}{c} | \\ \textcircled{\pi} \\ | \\ \textcircled{\pi} \\ | \\ \textcircled{\pi} \\ | \\ \textcircled{\pi} \\ | \\ \textcircled{\pi} \\ | \\ \textcircled{\pi} \end{array}$$
$$-2\sqrt{2}i \begin{array}{c} | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \circ \end{array} -2i \begin{array}{c} | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\frac{\pi}{2}} \\ | \\ \textcircled{\pi} \end{array}$$
$$+8\sqrt{2}i \begin{array}{c} | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \textcircled{\pi} \end{array} +8\sqrt{2}i \begin{array}{c} | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \square \\ | \\ \textcircled{\pi} \\ | \\ \square \\ | \\ \square \end{array}$$

Applying the decomposition

We pick some spiders to decompose and unfuse the phases:



And now we can apply the magic state decomposition.

Better decompositions

Improved upper bounds on the stabilizer rank of magic states

Hammam Qassim ^{*†‡ §}

Hakop Pashayan ^{*¶ ||}

David Gosset ^{*¶}

June 16, 2021

⇒ improved stabiliser decompositions, including 6-to-6 decomp (giving $\alpha \approx 0.431$ instead of $\alpha \approx 0.467$), and other decomps giving $\alpha < 0.40$.

Cat states

Qassim et al. uses *cat states*:

$$|\text{cat}_n\rangle := \frac{1}{\sqrt{2}}(\mathbb{I}^{\otimes n} + Z^{\otimes n})|T\rangle^{\otimes n} = \frac{1}{\sqrt{2^{n+1}}} \left(\textcircled{\frac{\pi}{4}}^{\otimes n} + \textcircled{\frac{5\pi}{4}}^{\otimes n} \right)$$

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These actually have nice representation in ZX:

$$|\text{cat}_n\rangle = \frac{1}{\sqrt{2}} \begin{array}{c} \textcircled{\frac{\pi}{4}} \\ \textcircled{\frac{\pi}{4}} \\ \vdots \\ \textcircled{\frac{\pi}{4}} \end{array}$$

Cat decompositions

They find good decomp of $|\text{cat}_6\rangle$:

The diagram shows the decomposition of the 6-qubit cat state $|\text{cat}_6\rangle$. On the left, a central white node is connected to six green nodes, each labeled $\frac{\pi}{4}$. On the right, the decomposition is shown as a sum of two terms. The first term is $\frac{1}{2}$ multiplied by a central green node labeled $-\frac{\pi}{2}$ connected to six white nodes. The second term is $\frac{ie^{i\pi/4}}{\sqrt{2}}$ multiplied by a central white node connected to six green nodes, each labeled $\frac{\pi}{2}$. The third term is $-\frac{e^{i\pi/4}}{\sqrt{2}}$ multiplied by a central white node connected to six green nodes, each labeled $\frac{\pi}{2}$.

$$|\text{cat}_6\rangle = \frac{1}{2} \left[\text{node} \left(-\frac{\pi}{2} \right) + \frac{ie^{i\pi/4}}{\sqrt{2}} \left[\text{node} \left(\frac{\pi}{2} \right) - \frac{e^{i\pi/4}}{\sqrt{2}} \left[\text{node} \left(\frac{\pi}{2} \right) \right] \right] \right]$$

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- A term with a coefficient $\frac{1}{2}$ and a central green node labeled $-\frac{\pi}{2}$ connected to six lines.
- A term with a coefficient $\frac{ie^{i\pi/4}}{\sqrt{2}}$ and a central white node connected to six green nodes, each labeled $\frac{\pi}{4}$.
- A term with a coefficient $-\frac{e^{i\pi/4}}{\sqrt{2}}$ and a central white node connected to six green nodes, each labeled $\frac{\pi}{2}$.

So, if we only have $|\text{cat}_6\rangle$ diagrams, we would have $\alpha \approx 0.264$.

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$$\begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ = \frac{1}{2} \begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ + \frac{ie^{i\pi/4}}{\sqrt{2}} \begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ - \frac{e^{i\pi/4}}{\sqrt{2}} \begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ$$

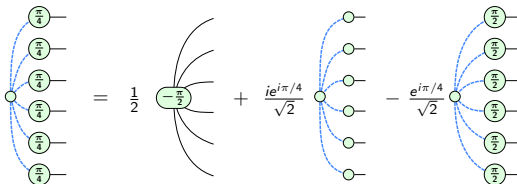
So, if we only have $|\text{cat}_6\rangle$ diagrams, we would have $\alpha \approx 0.264$.
But we can do even better!

$$\begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ = \frac{e^{-i\pi/4}}{\sqrt{2}} \begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ + i \begin{array}{c} \circ \\ \circ \\ \circ \\ \circ \end{array} \text{---} \circ$$

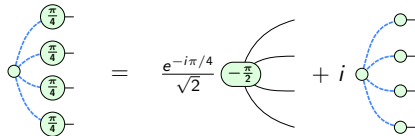
This would give $\alpha = 0.25$.

Cat decompositions

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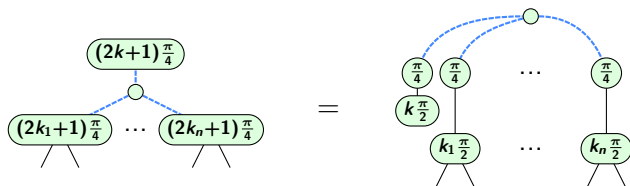


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Using these we get good decompositions for $|\text{cat}_k\rangle$ with $k \leq 6$.

Phase gadgets are cat states

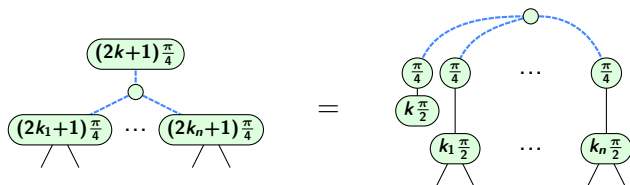
We find cat states as phase gadgets:



So n -legged phase gadget is $|\text{cat}_{n+1}\rangle$ state.

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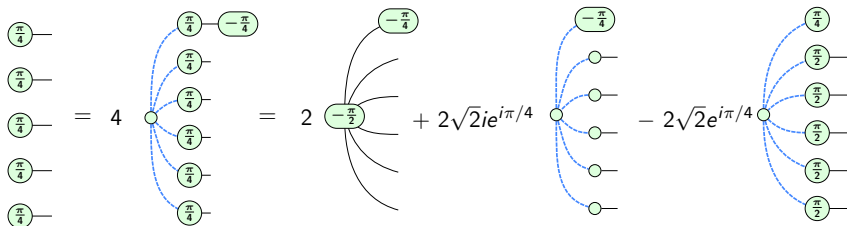
So n -legged phase gadget is $|\text{cat}_{n+1}\rangle$ state.

So as long as there are phase gadgets with ≤ 5 legs, we can use these decompositions.

Partial stabiliser decomp

But what if there are no phase gadgets?

Then we can do the following 'partial' decomp:



This trades 5 magic states for 3 terms with 1 magic per term.

So effectively removes 4 magic states.

This is then a 4-to-3 decomp: $\alpha \approx 0.396$.

Full strategy

We are hence looking for the following things to decompose:

1. a phase gadget with 3 legs ($\alpha = 0.25$),
2. a phase gadget with 5 legs ($\alpha \approx 0.264$),
3. a phase gadget with 4 legs ($\alpha \approx 0.317$),
4. a phase gadget with 2 legs ($\alpha = 1/3 \approx 0.333$),
5. any 5 T-spiders ($\alpha \approx 0.396$).

So how well does all this work?

Asymptotic benefit

- ▶ Worst case: no T-like phases killed during simplification.

Asymptotic benefit

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- ▶ (Bravyi et al., 2016) gave $O(2^{\alpha k} k^3)$.
- ▶ Benefit comes from preventing 'double work': we 'partially evaluate' the stabilisers by simplifying the diagrams.

Actual benefit

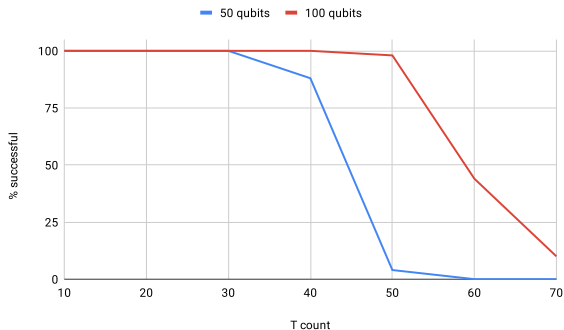
We benchmarked our method on two families of circuits:

- ▶ 50- and 100-qubit random Clifford+T circuits built out of Pauli exponentials.
- ▶ 50-qubit hidden-shift circuits (type of CCZ circuit).

We are sampling from the output distribution (using strong simulation).

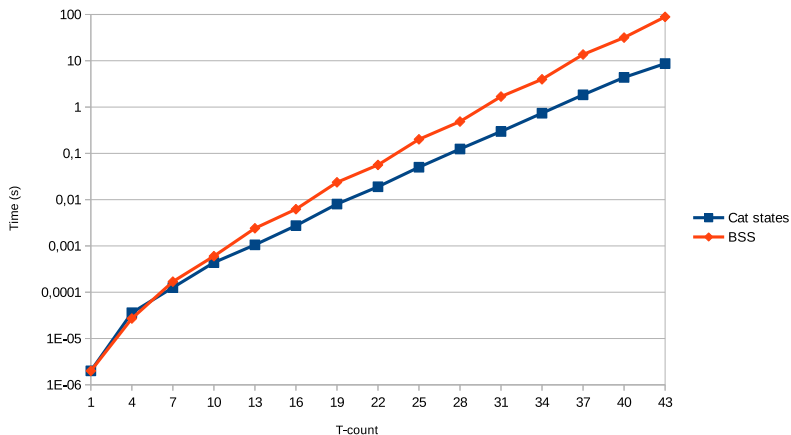
Code is implemented in `quixx`, a Rust port of PyZX.

Benchmark: Clifford+T



Percentage of random 50- and 100-qubit circuits of a given T-count that were successfully sampled in under 5 minutes. For each T-count 50 random circuits were generated.

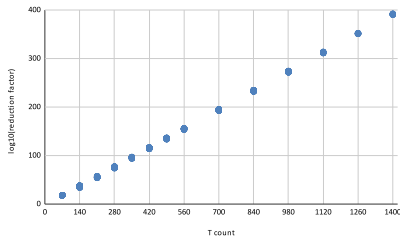
Benchmark: Clifford+T cat-decomp comparison



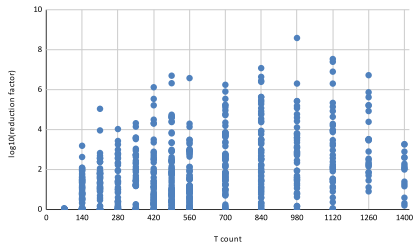
Runtime of random 20-qubit Clifford+T circuit simulations (avg of 10 runs per T-count).

Benchmark: hidden-shift circuit term reduction

Reduction vs. naïve BSS decomposition

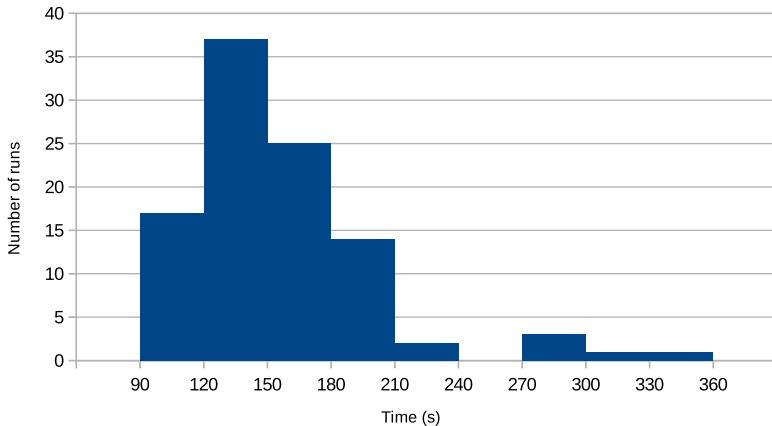


Reduction vs. simplified BSS decomposition



Reduction in term count on 50-qubit hidden shift circuits vs. naïve BSS decomposition (left) and BSS decomposition after single ZX-simplification (right).

Benchmark: hidden-shift 50-qubit simulation time



The time distribution of simulating 100 random 50-qubit hidden-shift circuits with T-count 1400 using our new decompositions.

Conclusions

- ▶ Using ZX we can greatly speed-up stabiliser rank simulations.
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Moral of the story: optimisation and simulation are not separate. They are two sides of the same coin.

Future work:

- ▶ Use more diagram optimisations and decompositions
- ▶ Find heuristics for picking good spiders to decompose.
- ▶ Approximate simulation and better weak simulation.
- ▶ Use *quantum measurement w/o computing marginals* technique.

Thank you for your attention!

Further reading:

- ▶ Kissinger & vdW. *Simulating quantum circuits with ZX-calculus reduced stabiliser decompositions.*
arXiv: 2109.01076
- ▶ Kissinger, Vilmart & vdW. *Classical simulation of quantum circuits with partial and graphical stabiliser decompositions.*
arXiv: 2202.09202
- ▶ Qassim, Pashayan, Gosset. *Improved upper bounds on the stabilizer rank of magic states.*
arXiv: 2106.07740
- ▶ Bravyi, Gosset. *Improved classical simulation of quantum circuits dominated by Clifford gates.*
arXiv: 1601.07601