

Supercharging classical simulation of quantum circuits with the ZX-calculus

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Quantum circuit simulation

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 - ▶ contraction order finding
 - ▶ edge cutting
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 - ▶ ...
- ▶ Stabiliser-decomposition based methods:
 - ▶ Stabiliser extent
 - ▶ Stabiliser rank ← **This talk**

Stabiliser rank

Gottesman-Knill Theorem

A quantum computation consisting of

- ▶ stabiliser states,
- ▶ Clifford unitaries,
- ▶ stabiliser measurements,

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Observation: Stabiliser states span entire space of states. So:

$$|\psi\rangle = \sum_{j=1}^{\chi} \lambda_j |\psi_j\rangle \text{ where } |\psi_j\rangle \text{ is stabiliser.}$$

χ is called **stabiliser rank** of $|\psi\rangle$.

Simulating using stabiliser rank

- ▶ Start with Clifford+T circuit.
- ▶ Write each T gate as magic state injection.
- ▶ Decompose T states into sum of stabilisers.
- ▶ Efficiently simulate resulting Clifford circuits.
- ▶ Add results together.
- ▶ We're done!

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What's the catch?

Stabiliser rank of k T states scales exponentially with k .

... But it's not just 2^k terms. We can do better!

Stabiliser ranks of T magic states

Recall $|T\rangle \propto |0\rangle + e^{i\pi/4}|1\rangle$.

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Turns out $\chi(|T\rangle^{\otimes 6}) \leq 7$ so

$\chi(|T\rangle^{\otimes k}) \leq 2^{\alpha k}$ where $\alpha = \log_2(7)/6 \approx 0.467$.

Found by Bravyi, Smith, Smolin (BSS) in 2016.

Our idea

Do this simulation, but with *ZX-diagrams* instead of circuits!

Benefit: optimise intermediate ZX-diagrams to reduce T-count.

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Benefit: optimise intermediate ZX-diagrams to reduce T-count.

Described in two papers:

- ▶ Kissinger & vdW. *Simulating quantum circuits with ZX-calculus reduced stabiliser decompositions.*
arXiv:2109.01076
- ▶ Kissinger, Vilmart & vdW. *Classical simulation of quantum circuits with partial and graphical stabiliser decompositions.*
arXiv:2202.09202

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Let's go through these steps in more detail.

Writing Clifford+T circuit as ZX-diagram

$$\text{CNOT} = \sqrt{2} \begin{array}{c} \text{---} \text{---} \text{---} \text{---} \\ \text{---} \text{---} \text{---} \text{---} \end{array} \begin{array}{c} \text{---} \text{---} \text{---} \text{---} \\ \text{---} \text{---} \text{---} \text{---} \end{array} \begin{array}{c} \text{---} \text{---} \text{---} \text{---} \\ \text{---} \text{---} \text{---} \text{---} \end{array} \quad \text{S} = \text{---} \text{---} \text{---} \text{---} \quad \text{Had} = \text{---} \text{---} \text{---} \text{---}$$

$$\text{T} = \text{---} \text{---} \text{---} \text{---} \quad |x\rangle = \frac{1}{\sqrt{2}} \text{---} \text{---} \text{---} \text{---} \quad \langle x| = \frac{1}{\sqrt{2}} \text{---} \text{---} \text{---} \text{---} \quad (x \in \{0, 1\})$$

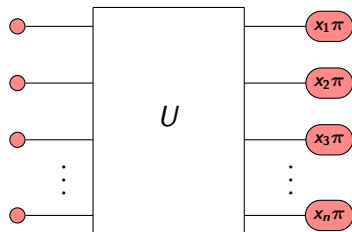
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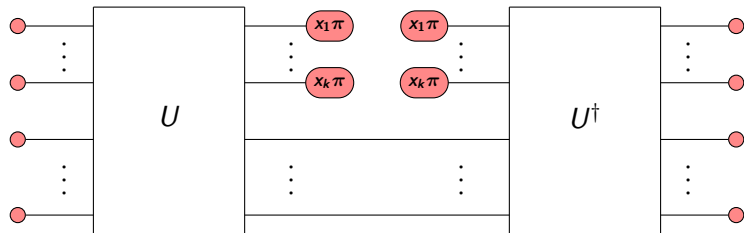
Calculating single amplitude:

$$\langle \vec{x} | U | 0 \dots 0 \rangle = \left(\frac{1}{\sqrt{2}} \right)^{2n}$$



Marginal probabilities

To calculate marginal probability, use *doubling* technique:



Strong simulation vs Weak simulation

Weak sim: approx. sample from the same output distribution.

Strong sim: approx. calculate any marginal probability.

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Strong sim: approx. calculate any marginal probability.

We are doing *exact* strong simulation here.

- ▶ Writing circuit as ZX-diagram ✓
- ▶ Optimising ZX-diagram
- ▶ Decomposing magic states

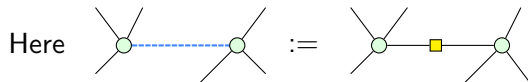
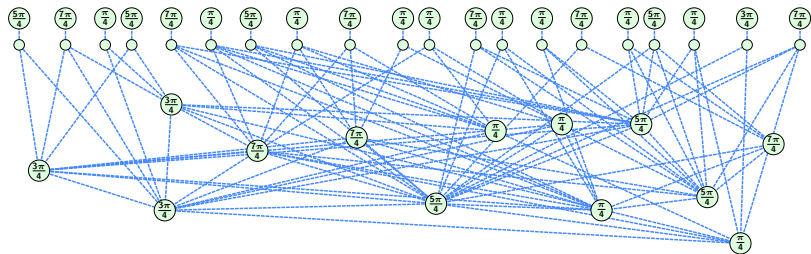
Simplifying ZX-diagrams

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Reducing T-count with the ZX-calculus.

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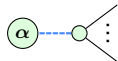
This transforms every ZX-diagram into something like this:



Properties of reduced diagram

The important part:

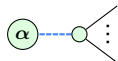
- ▶ Every spider carries a non-Clifford phase,
- ▶ or is part of a *phase gadget*:



Properties of reduced diagram

The important part:

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Particularly: if original circuit had k T gates, resulting diagram has $\leq 2k$ spiders (regardless of #qubits or #gates).

- ▶ Writing circuit as ZX-diagram ✓
- ▶ Optimising ZX-diagram ✓
- ▶ Decomposing magic states

Decomposing T-like spiders

The 6-to-7 magic state decomposition in ZX is:

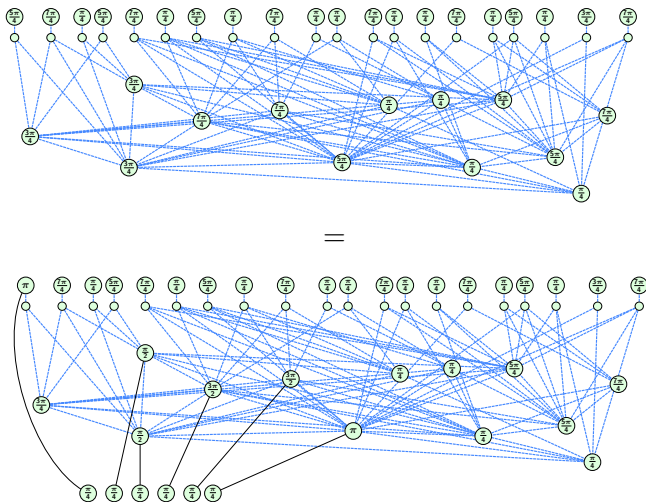
$$\begin{aligned}
 e^{i\pi/4} \begin{array}{c} | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \\ | \\ \bigcirc(\frac{\pi}{4}) \end{array} &= +2e^{i\pi/4} \begin{array}{c} | \quad | \quad | \quad | \quad | \quad | \\ \bigcirc(\frac{\pi}{2}) \end{array} \\
 &- \frac{1+\sqrt{2}}{4} \begin{array}{c} | \\ \bigcirc \\ | \\ \bigcirc \\ | \\ \bigcirc \\ | \\ \bigcirc \\ | \\ \bigcirc \end{array} + \frac{1-\sqrt{2}}{4} \begin{array}{c} | \\ \bigcirc(\pi) \\ | \\ \bigcirc(\pi) \\ | \\ \bigcirc(\pi) \\ | \\ \bigcirc(\pi) \\ | \\ \bigcirc(\pi) \end{array} \\
 &- 2\sqrt{2}i \begin{array}{c} | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \end{array} - 2i \begin{array}{c} | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \\ | \\ \bigcirc(\frac{\pi}{2}) \end{array} \\
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 \end{aligned}$$

The diagram illustrates the decomposition of a 6-to-7 T-like spider into several terms. The leftmost spider has six legs, each ending in a circle containing $\pi/4$. The rightmost spider has six legs, each ending in a circle containing $\pi/2$. The decomposition terms are:

- $2e^{i\pi/4}$ times a spider with six legs and a central circle containing $\pi/2$.
- $-\frac{1+\sqrt{2}}{4}$ times a spider with six legs and six central circles, each containing $\pi/4$.
- $+\frac{1-\sqrt{2}}{4}$ times a spider with six legs and six central circles, each containing π .
- $-2\sqrt{2}i$ times a spider with six legs and six central circles, each containing $\pi/2$, with blue dashed arcs connecting the circles.
- $-2i$ times a spider with six legs and six central circles, each containing $\pi/2$, with blue dashed arcs connecting the circles.
- $+8\sqrt{2}i$ times a spider with six legs and six central circles, each containing $\pi/4$, with blue dashed arcs connecting the circles.
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Applying the decomposition

We pick some spiders to decompose and unfuse the phases:



And now we can apply the magic state decomposition.

Full algorithm

- ▶ After first decomposition we have 7 diagrams.
- ▶ Repeat whole procedure for each of these diagrams, end up with $7 * 7$ diagrams.
- ▶ Repeat until there aren't 6 magic states to decompose.
- ▶ Diagrams are then small enough to calculate value of directly.
- ▶ Sum values to get the amplitude/prob we were calculating.

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But we can do better...

Better decompositions

Improved upper bounds on the stabilizer rank of magic states

Hammam Qassim ^{*†‡} §

Hakop Pashayan ^{*¶} ||

David Gosset ^{*¶}

June 16, 2021

⇒ improved stabiliser decompositions, including 6-to-6 decomp (giving $\alpha \approx 0.431$ instead of $\alpha \approx 0.467$), and other decomp giving $\alpha < 0.40$.

Cat states

Qassim et al. uses *cat states*:

$$|\text{cat}_n\rangle := \frac{1}{\sqrt{2}}(\mathbb{I}^{\otimes n} + Z^{\otimes n})|T\rangle^{\otimes n} = \frac{1}{\sqrt{2^{n+1}}} \left(\textcircled{\frac{\pi}{4}}^{\otimes n} + \textcircled{\frac{5\pi}{4}}^{\otimes n} \right)$$

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These actually have nice representation in ZX:

$$|\text{cat}_n\rangle = \frac{1}{\sqrt{2}} \begin{array}{c} \textcircled{\frac{\pi}{4}} \\ \textcircled{\frac{\pi}{4}} \\ \vdots \\ \textcircled{\frac{\pi}{4}} \end{array}$$

Cat decompositions

They find good decomp of $|\text{cat}_6\rangle$:

The diagram shows the decomposition of the 6-qubit cat state $|\text{cat}_6\rangle$. On the left, a central white node is connected to six green nodes, each labeled $\frac{\pi}{4}$. This is equal to the sum of two terms. The first term is $\frac{1}{2}$ times a central green node labeled $-\frac{\pi}{2}$ connected to six white nodes. The second term is $\frac{ie^{i\pi/4}}{\sqrt{2}}$ times a central white node connected to six green nodes, each labeled $\frac{\pi}{2}$. The third term is $-\frac{e^{i\pi/4}}{\sqrt{2}}$ times a central white node connected to six green nodes, each labeled $\frac{\pi}{2}$.

$$\begin{array}{c} \frac{\pi}{4} \\ \frac{\pi}{4} \\ \frac{\pi}{4} \\ \frac{\pi}{4} \\ \frac{\pi}{4} \\ \frac{\pi}{4} \end{array} = \frac{1}{2} \begin{array}{c} -\frac{\pi}{2} \end{array} + \frac{ie^{i\pi/4}}{\sqrt{2}} \begin{array}{c} \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \end{array} - \frac{e^{i\pi/4}}{\sqrt{2}} \begin{array}{c} \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \\ \frac{\pi}{2} \end{array}$$

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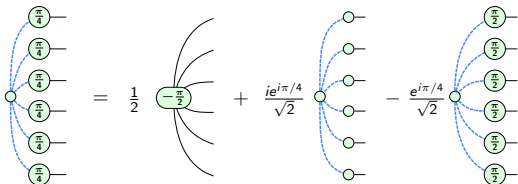
The diagram shows the decomposition of the 6-qubit cat state $|\text{cat}_6\rangle$. On the left, a central white node is connected to six green nodes, each labeled $\frac{\pi}{4}$. This is equal to the sum of three terms:

- A term with a coefficient $\frac{1}{2}$ and a central green node labeled $-\frac{\pi}{2}$ connected to six lines.
- A term with a coefficient $\frac{ie^{i\pi/4}}{\sqrt{2}}$ and a central white node connected to six green nodes, each labeled $\frac{\pi}{4}$.
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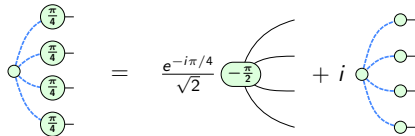
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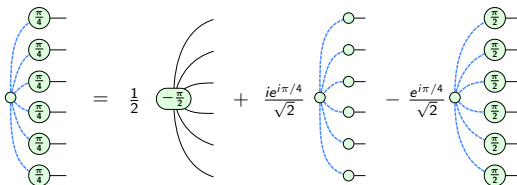
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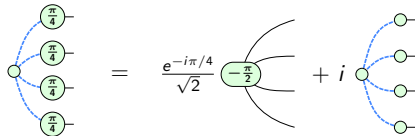
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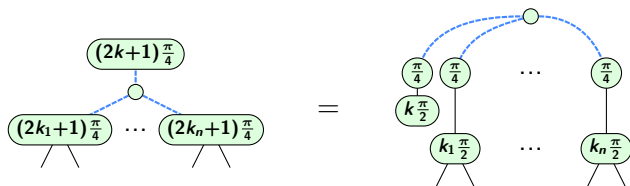


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Using these we get good decompositions for $|\text{cat}_k\rangle$ with $k \leq 6$.

Phase gadgets are cat states

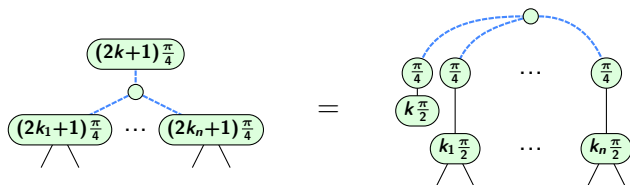
We find cat states as phase gadgets:



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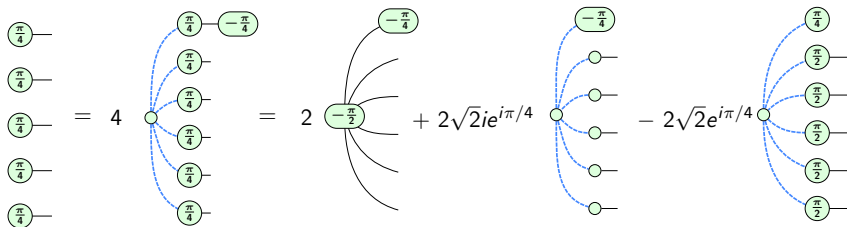
So n -legged phase gadget is $|\text{cat}_{n+1}\rangle$ state.

So as long as there are phase gadgets with ≤ 5 legs, we can use these decompositions.

Partial stabiliser decomp

But what if there are no phase gadgets?

Then we can do the following 'partial' decomp:



This trades 5 magic states for 3 terms with 1 magic per term.

So effectively removes 4 magic states.

This is then a 4-to-3 decomp: $\alpha \approx 0.396$.

Full strategy

We are hence looking for the following things to decompose:

1. a phase gadget with 3 legs ($\alpha = 0.25$),
2. a phase gadget with 5 legs ($\alpha \approx 0.264$),
3. a phase gadget with 4 legs ($\alpha \approx 0.317$),
4. a phase gadget with 2 legs ($\alpha = 1/3 \approx 0.333$),
5. any 5 T-spiders ($\alpha \approx 0.396$).

So how well does all this work?

Asymptotic benefit

- ▶ Worst case: no T-like phases killed during simplification.

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- ▶ There are $O(2^{\alpha k})$ diagrams, so total cost is $O(2^{\alpha k} k^2)$.
- ▶ (Bravyi et al., 2016) gave $O(2^{\alpha k} k^3)$.
- ▶ Benefit comes from preventing 'double work': we 'partially evaluate' the stabilisers by simplifying the diagrams.

Actual benefit

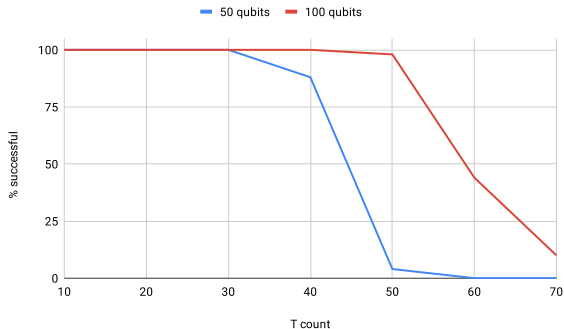
We benchmarked our method on two families of circuits:

- ▶ 50- and 100-qubit random Clifford+T circuits built out of Pauli exponentials.
- ▶ 50-qubit hidden-shift circuits (type of CCZ circuit).

We are sampling from the output distribution (using strong simulation).

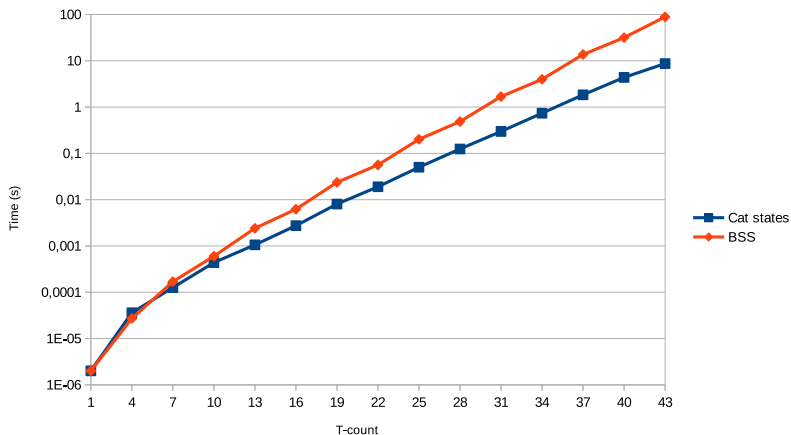
Code is implemented in `quixx`, a Rust port of PyZX.

Benchmark: Clifford+T



Percentage of random 50- and 100-qubit circuits of a given T-count that were successfully sampled in under 5 minutes. For each T-count 50 random circuits were generated.

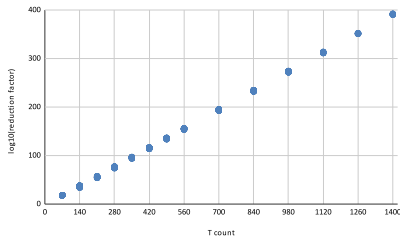
Benchmark: Clifford+T cat-decomp comparison



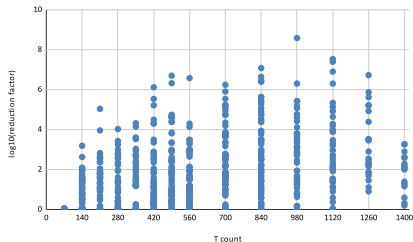
Runtime of random 20-qubit Clifford+T circuit simulations (avg of 10 runs per T-count).

Benchmark: hidden-shift circuit term reduction

Reduction vs. naïve BSS decomposition

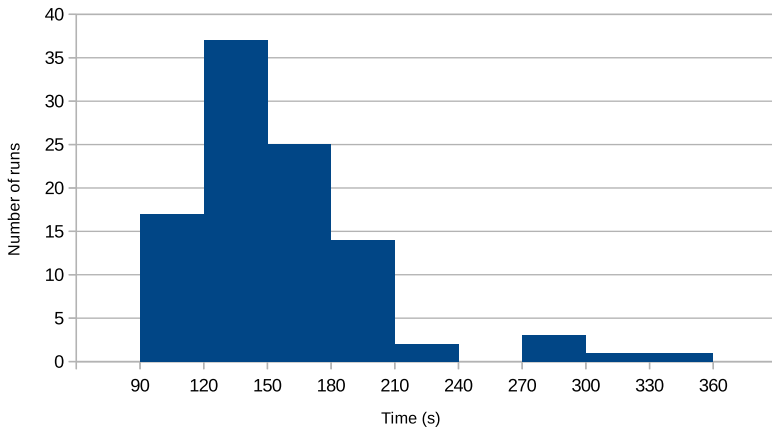


Reduction vs. simplified BSS decomposition



Reduction in term count on 50-qubit hidden shift circuits vs. naïve BSS decomposition (left) and BSS decomposition after single ZX-simplification (right).

Benchmark: hidden-shift 50-qubit simulation time



The time distribution of simulating 100 random 50-qubit hidden-shift circuits with T-count 1400 using our new decompositions.

Conclusions

- ▶ Using ZX we can greatly speed-up stabiliser rank simulations.
- ▶ Especially for structured circuits.
- ▶ It allows us to use better decompositions for substructures of diagrams, and to introduce partial stabiliser decompositions.

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- ▶ Especially for structured circuits.
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Moral of the story: optimisation and simulation are not separate.
They are two sides of the same coin.

Conclusions

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Moral of the story: optimisation and simulation are not separate. They are two sides of the same coin.

Future work:

- ▶ Use more diagram optimisations.
- ▶ Find heuristics for picking good spiders to decompose.
- ▶ Approximate simulation and better weak simulation.
- ▶ stabiliser *extent* methods.
- ▶ Use *quantum measurement w/o computing marginals* technique.

Thank you for your attention!

Further reading:

- ▶ Kissinger & vdW. *Simulating quantum circuits with ZX-calculus reduced stabiliser decompositions.*
arXiv: 2109.01076
- ▶ Kissinger, Vilmart & vdW. *Classical simulation of quantum circuits with partial and graphical stabiliser decompositions.*
arXiv: 2202.09202
- ▶ Qassim, Pashayan, Gosset. *Improved upper bounds on the stabilizer rank of magic states.*
arXiv: 2106.07740
- ▶ Bravyi, Gosset. *Improved classical simulation of quantum circuits dominated by Clifford gates.*
arXiv: 1601.07601