

Dichotomy between deterministic and probabilistic models in countably additive effectus theory

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Generalized Probabilistic Theories

GPTs are generalisations of quantum theory.

They consist of

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- ▶ the 'empty system' I ,
- ▶ operations $f : A \rightarrow B$,

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- ▶ and the **scalars** $\{s : I \rightarrow I\}$ are the real unit interval $[0, 1]$

Special operations:

- ▶ **States** $\text{St}(A) := \{\omega : I \rightarrow A\}$
- ▶ **Effects** $\text{Eff}(A) := \{p : A \rightarrow I\}$
- ▶ $p \circ \omega$ is probability that p holds on state ω

From GPTs to effectuses

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Solution: allow more general sets of scalars $\{s : I \rightarrow I\}$.

Result: effectus theory.

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Definition

A **partial commutative monoid** (PCM) $(X, \otimes, 0)$ is a set X with a *partial* associative commutative operation \otimes with unit 0 .

$$(x \otimes y) \otimes z = x \otimes (y \otimes z) \quad x \otimes y = y \otimes x \quad x \otimes 0 = x$$

Write $x \perp y$ when $x \otimes y$ is defined.

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- ▶ Let **Cstar** be category of unital C^* -algebras with positive subunital maps.
Homsets **Cstar**($\mathfrak{A}, \mathfrak{B}$) are PCMs:
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Pfn and **Cstar** are examples of **PCM-enriched categories**:

$$(f \vee g) \circ h = (f \circ h) \vee (g \circ h) \quad h \circ (f \vee g) = (h \circ f) \vee (h \circ g)$$

Effect algebras

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Examples

- ▶ $[0, 1]$ with $a^\perp := 1 - a$.
- ▶ A Boolean algebra: $a \perp b$ when $a \wedge b = 0$ and then $a \oplus b = a \vee b$. a^\perp is the regular negation.
- ▶ $\mathbf{Cstar}(\mathbb{C}, \mathfrak{A}) \cong [0, 1]_{\mathfrak{A}}$ with $a^\perp := 1 - a$.

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Examples:

- ▶ **Pfn**: $I = \{*\}$, $\mathbf{Pfn}(A, I) \cong \mathcal{P}(A)$.
- ▶ **Cstar^{op}**: $I = \mathbb{C}$, $\mathbf{Cstar}^{\text{op}}(\mathfrak{A}, I) \cong [0, 1]_{\mathfrak{A}}$.

Relating effectus to GPTs

Let \mathbf{C} be an effectus, and $A \in \mathbf{C}$.

- ▶ Effects $\text{Eff}(A) := \mathbf{C}(A, I)$ form effect algebra.
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- ▶ What are they in an effectus?
- ▶ $\mathbf{Pfn}(I, I) \cong \{0, 1\}$ & $\mathbf{Cstar}^{\text{op}}(I, I) \cong [0, 1]$.
- ▶ In general: $\mathbf{C}(I, I)$ is effect algebra.
- ▶ But also has a 'multiplication' given by composition $I \xrightarrow{s} I \xrightarrow{t} I$.

Scalars in effectus

Definition

An **effect monoid** $(M, \otimes, 0, 1, \cdot)$ is an effect algebra with associative distributive multiplication:

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Examples:

- ▶ $[0, 1]$.
- ▶ Any Boolean algebra: $a \otimes b := a \vee b$, $a \cdot b := a \wedge b$.
- ▶ $\{f : X \rightarrow [0, 1] \text{ continuous}\}$ for a compact Hausdorff space X (i.e. unit interval of commutative unital C^* -algebra).

Effectus over an effect monoid

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- ▶ An M -**effect module** E is an effect algebra with suitable left M -action $\cdot : M \times E \rightarrow E$.
- ▶ $\mathbf{EMod}_M^{\text{op}}$ is an effectus. Eff: $\mathbf{C} \rightarrow \mathbf{EMod}_M^{\text{op}}$ is a functor.

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- ▶ $\text{St}(A)$ has 'weight function' $|\omega| := \mathbf{1}_A \circ \omega$.
- ▶ A **weight M -module** X is a PCM with M -action $\cdot : M \times X \rightarrow X$ and suitable *weight* function $|\cdot| : X \rightarrow M$.
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Any effect monoid is the set of scalars of some effectus

Corollary: There are some weird effectuses out there

σ -PCMs

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Definition (informal)

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Definition (informal)

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Examples:

- ▶ **Pfn**(A, B).
- ▶ Let **Wstar** be category of von Neumann algebras with normal positive subunital maps. Then **Wstar**($\mathfrak{A}, \mathfrak{B}$) is σ -PCM.

In fact: **Pfn** and **Wstar** are σ -**PCM** enriched.

σ -effectus

Definition

σ -effectus is σ -PCM-enriched effectus with countable coproducts.

Examples:

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It turns out that σ -effectuses are way more well-behaved.

σ -effect monoids

Proposition

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Theorem (Westerbaan, Westerbaan & vdW, LICS'20)

An ω -directed-complete effect monoid M embeds into $M_1 \oplus M_2$
where M_1 is a ω -complete Boolean algebra and
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Corollary

Scalars in a σ -effectus are commutative.

Normalisation in σ -effectuses

Theorem

Let \mathbf{C} be a σ -effectus with $M = \mathbf{C}(I, I)$.

The following are equivalent.

- ▶ States in \mathbf{C} can be normalized.
- ▶ Non-zero scalars are epi.
- ▶ M has a 'division' operation.
- ▶ M has no zero divisors ($a \cdot b = 0 \implies a = 0$ or $b = 0$).
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Furthermore, if any and thus all these conditions hold then

$M \cong \{0\}$, $M \cong \{0, 1\}$ or $M \cong [0, 1]$.

Dichotomy between deterministic and probabilistic models

Hence: σ -effectuses with normalization come in three types:

- ▶ $\mathbf{C}(I, I) \cong \{0\}$: only holds when \mathbf{C} is equivalent to the trivial single-object category with a single morphism.

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When appropriate operational *separation properties* are satisfied, we can say even more about these latter two cases.

Separation properties

Definition

An effectus \mathbf{C} has **state-separation** when $f \circ \omega = g \circ \omega$ for all states $\omega \in \text{Eff}(A)$ implies $f = g$ for any $f, g : A \rightarrow B$.
(i.e. iff $\text{St} : \mathbf{C} \rightarrow \mathbf{WMod}_M$ faithful)

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Non-Example

Category of effect algebras $\mathbf{EA}^{\text{op}} \cong \mathbf{EMod}_{\{0,1\}}^{\text{op}}$.

Let $P(\mathcal{H}) \in \mathbf{EA}$ be the projections on a Hilbert space \mathcal{H} with $\dim(\mathcal{H}) > 2$. Then Kochen-Specker theorem says

$\text{St}(P(\mathcal{H})) = \{0\}$.

Classical deterministic effectuses

Theorem

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Let \mathbf{C} be a σ -effectus with $\mathbf{C}(I, I) \cong \{0, 1\}$ and state-separation. Then there is a faithful morphism of σ -effectuses $F : \mathbf{C} \rightarrow \mathbf{Pfn}$, and $\text{St}(A) \cong F(A)$ for all $A \in \mathbf{C}$.

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Hence, such effectuses are entirely classical.

Corollary

Nonclassical σ -effectuses w/ normalisation **must** have scalars $[0, 1]$.

Convex embeddings of effectuses

A similar sort of embedding holds in the probabilistic setting.

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A **Banach** OUS is furthermore complete in its canonical norm.

A σ -OUS furthermore has countable directed suprema.

Convex embeddings of effectuses

A similar sort of embedding holds in the probabilistic setting.

Definition (informal)

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Theorem

Let \mathbf{C} be a σ -effectus with $\mathbf{C}(I, I) \cong [0, 1]$ and effect-separation.

Then there is a faithful morphism of σ -effectuses

$F : \mathbf{C} \rightarrow \sigma\mathbf{BOUS}^{\text{op}}$, the category of Banach σ -OUSes.

Hence: Probabilistic effectuses embed into the well-studied land of real ordered vector spaces.

Summary

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Summary

- ▶ A non-trivial σ -effectus with normalisation is either deterministic or probabilistic.
- ▶ When it is deterministic and has state-separation then it embeds into **Pfn** and **BA**^{op}, and hence is classical.
- ▶ When it is probabilistic and has effect-separation it embeds into **BOUS**^{op}, and hence reduces to a standard GPT-like object.
- ▶ So we managed to go from an abstract categorical framework to concrete well-studied standard settings.

Future work

- ▶ Study general σ -effectuses (without normalization).
- ▶ Find inherent categorical characterisation of σ -effectuses.

Thank you for your attention

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