# Quantum Compilation using the ZX-calculus 

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## The ZX-calculus

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Used in:

- Quantum circuit optimisation and compilation
- Measurement-based quantum computation
- Surface codes and lattice surgery


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It is also a convenient tool for day-to-day quantum reasoning

## Further Reading

For references and details see:
ZX-calculus for the working quantum computer scientist
https://arxiv.org/abs/2012.13966

For a book-length introduction see:
Picturing Quantum Processes: A First Course in Quantum Theory and Diagrammatic Reasoning
by Bob Coecke and Aleks Kissinger

## Quantum circuits



## Circuit identities


$\sqrt{\mathrm{H}} \sqrt{\mathrm{H}}=$
$\square$
$\sqrt{\mathrm{T}^{+}}-\sqrt{\mathrm{T}}=$
$-\sqrt{T}-\sqrt{T}=-\sqrt{S}$

## Gate commutation



## More circuit equalities

$$
\begin{aligned}
& \because=\text { 二 } \\
& \text { - } \\
& \text { - } \ddagger=\text { ! } \ddagger
\end{aligned}
$$

$$
\begin{aligned}
& \text { ! } \sqrt{H}!=\overline{\sqrt{S}-(\mathbb{H} \cdot \sqrt{S}-\sqrt{S-T-S}} \cdot \omega^{-1}
\end{aligned}
$$

*Selinger 2015

## And more circuit equalities

$$
\begin{aligned}
& - \text { 四 }- \text { 困 }-=- \text { 困 } \\
& \text {-四困- = - 困- } \\
& - \text { 四- } \\
& \text { ? 四- - 四! }
\end{aligned}
$$

$$
\begin{aligned}
& \text { • }
\end{aligned}
$$

$$
\begin{aligned}
& \text { - }
\end{aligned}
$$

$$
\begin{aligned}
& \text { - } \mathbb{⿴ 囗 ⿻}
\end{aligned}
$$

$$
\begin{aligned}
& - \text { 표 }
\end{aligned}
$$

$$
\begin{aligned}
& - \text { 因國 }- \text { - } \\
& - \text { 因國 }=- \text { - }
\end{aligned}
$$







F




 F


＊Selinger 2015

## And even more circuit equalities

$$
\begin{aligned}
& R_{7}: \sqrt{T^{8}}=-\frac{\sqrt{U^{4}}}{-\sqrt{V^{4}}-}=\frac{-\sqrt{T^{4}}}{-\sqrt{T^{4}}} \\
& R_{10}: \omega^{8}=\quad R_{11}:-x-T-x-2-\sqrt{Y^{7}}-
\end{aligned}
$$

$$
\begin{aligned}
& R_{12}: \frac{1}{x} \frac{1}{x}=\sqrt{T}
\end{aligned}
$$

*Amy, Chen, \& Ross 2018

## Quantum circuits bad!

Why is this so terrible?

- Choice of gates is a bit arbitrary.
- The notation is not "quantum native".
- Wires are rigid going from left-to-right.


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Why is this so terrible?

- Choice of gates is a bit arbitrary.
- The notation is not "quantum native".
- Wires are rigid going from left-to-right.

The ZX-calculus essentially gets rid of these problems.

## ZX-diagrams

On a surface level, ZX-diagrams are alternative notation to circuits


## Circuit identity in ZX



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dots of same colour commute through each other

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dots of same colour commute through each other

$\leadsto$


## Circuit identity in ZX


dots of same colour commute through each other

$\leadsto$


More fundamental rule: dots of same colour fuse


States in ZX

$$
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|+\rangle=\frac{1}{\sqrt{2}}\binom{1}{1} \quad|-\rangle=\frac{1}{\sqrt{2}}\binom{1}{-1}
$$

States in ZX

$$
\begin{gathered}
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|+\rangle=\frac{1}{\sqrt{2}}\binom{1}{1} \quad|-\rangle=\frac{1}{\sqrt{2}}\binom{1}{-1} \\
|0\rangle-\quad=|0\rangle \\
=-\quad
\end{gathered}
$$

States in ZX

$$
\begin{gathered}
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|+\rangle=\frac{1}{\sqrt{2}}\binom{1}{1} \quad|-\rangle=\frac{1}{\sqrt{2}}\binom{1}{-1} \\
|0\rangle-\quad=|0\rangle- \\
-\quad \\
|0\rangle-\quad \cdots \infty
\end{gathered}
$$

States in ZX

$$
\begin{gathered}
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|+\rangle=\frac{1}{\sqrt{2}}\binom{1}{1} \quad|-\rangle=\frac{1}{\sqrt{2}}\binom{1}{-1} \\
|0\rangle-\quad=|0\rangle- \\
-\quad-\quad \mid
\end{gathered}
$$

single-wire dot copies through opposite-coloured dot


$$
=\underline{Q^{o-}}=\square^{0-}=\underline{0^{o-}}
$$

States in ZX

$$
\begin{gathered}
|0\rangle=\binom{1}{0} \quad|1\rangle=\binom{0}{1} \quad|+\rangle=\frac{1}{\sqrt{2}}\binom{1}{1} \quad|-\rangle=\frac{1}{\sqrt{2}}\binom{1}{-1} \\
|0\rangle-\frac{0}{-\infty}=|0\rangle- \\
|0\rangle-
\end{gathered}
$$

single-wire dot copies through opposite-coloured dot

all rules hold with colours interchanged


Now let's formally introduce ZX-diagrams

## Spiders

What gates are to circuits, spiders are to ZX-diagrams.

## Spiders

What gates are to circuits, spiders are to ZX-diagrams.

$$
\begin{array}{cc}
\text { Z-spider } & \text { X-spider } \\
|0 \cdots 0\rangle\langle 0 \cdots 0| & |+\cdots+\rangle\langle+\cdots+| \\
+e^{i \alpha}|1 \cdots 1\rangle\langle 1 \cdots 1| & +e^{i \alpha}|-\cdots-\rangle\langle-\cdots-|
\end{array}
$$

## Spiders

What gates are to circuits, spiders are to ZX-diagrams.

$$
\begin{gathered}
\text { Z-spider } \\
|0 \cdots 0\rangle\langle 0 \cdots 0| \\
+e^{i \alpha}|1 \cdots 1\rangle\langle 1 \cdots 1| \\
\vdots \alpha
\end{gathered}
$$

X-spider
$|+\cdots+\rangle\langle+\cdots+|$
$+e^{i \alpha}|-\cdots-\rangle\langle-\cdots \cdot|$


For example:
$—$ —— $=|0\rangle\langle 0|+e^{i \alpha}|1\rangle\langle 1|=\left(\begin{array}{ll}1 & 0 \\ 0 & 0\end{array}\right)+\left(\begin{array}{cc}0 & 0 \\ 0 & e^{i \alpha}\end{array}\right)=\left(\begin{array}{cc}1 & 0 \\ 0 & e^{i \alpha}\end{array}\right)$
$-(\alpha)=|+\rangle\langle+|+e^{i \alpha}|-\rangle\langle-|=\frac{1}{2}\left(\begin{array}{ll}1 & 1 \\ 1 & 1\end{array}\right)+\frac{1}{2} e^{i \alpha}\left(\begin{array}{cc}1 & -1 \\ -1 & 1\end{array}\right)$

## Spiders cont.

If $\alpha=0$ we drop the label:



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If $\alpha=0$ we drop the label:


Example:

We ignore these non-zero scalar factors

## Formal composition

Spiders can be composed in two ways.

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Spiders can be composed in two ways. Horizontal composition gives tensor product:

$$
\begin{aligned}
& =\left(\begin{array}{ll}
1 & 0 \\
0 & 0 \\
0 & 0 \\
0 & 1
\end{array}\right) \\
& \square=\left(\begin{array}{ll}
1 & 0 \\
0 & 0 \\
0 & 0 \\
0 & 1
\end{array}\right) \otimes\left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right)=\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right)
\end{aligned}
$$

## Formal composition

Other tensor product:

$$
\begin{aligned}
\overline{0} & =\left(\begin{array}{ll}
1 & 0 \\
0 & 1
\end{array}\right) \otimes \frac{1}{\sqrt{2}}\left(\begin{array}{llll}
1 & 0 & 0 & 1 \\
0 & 1 & 1 & 0
\end{array}\right) \\
& =\frac{1}{\sqrt{2}}\left(\begin{array}{llllllll}
1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0
\end{array}\right)
\end{aligned}
$$

## Formal composition

Horizontal composition is regular composition of linear maps:


$$
\frac{1}{\sqrt{2}}\left(\begin{array}{llllllll}
1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0
\end{array}\right)\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1
\end{array}\right)=\frac{1}{\sqrt{2}}\left(\begin{array}{cccc}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right)
$$

## Building ZX-diagrams

Any ZX-diagram is built by simply iterating these vertical and horizontal compositions

## Symmetries

Note:


Hence, we may write


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Hence, we may write


In general: only connectivity matters


## ZX-diagrams summary

- Two types of generators: Z-spiders and X-spiders
- Can compose both horizontally and vertically
- Wires can connect every which way


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How powerful are ZX-diagrams as a representation?
Theorem
ZX-diagrams are universal: any linear map between qubits can be represented as a ZX-diagram.

So far it's just notation. What can we do with it?

## Rules for ZX-diagrams: The ZX-calculus



$$
-\mathrm{O}=
$$

$$
-\square \square-\quad=\quad
$$

$\alpha, \beta \in[0,2 \pi]$

## Spider fusion



Connected spiders of same colour fuse

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$$
-(\beta-\beta=-\alpha+\beta
$$

## State and pi-copy


$\pi$ 's and states copy through the other colour

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$\pi$ 's and states copy through the other colour
Combining rules:


## Hadamards and colour-changing

Definition of Hadamard in ZX:

$$
\square-:=-\left(\frac{\pi}{2}\right)-\frac{\pi}{2}-\frac{\pi}{2}-
$$

## Hadamards and colour-changing

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Rules:


## Hadamards and colour-changing

Definition of Hadamard in ZX:

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$$

Rules:


Derived rule: commuting Hadamards changes colour


## Hadamards and colour-changing

Definition of Hadamard in ZX:

$$
\square-\square:=-\frac{\pi}{2}-\frac{\pi}{2}-\frac{\pi}{2}-
$$

Rules:


Derived rule: commuting Hadamards changes colour


Consequence: Everything in ZX holds with colours reversed

## Bialgebra

Z-spiders act like COPY; X-spiders act like XOR:


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Z-spiders act like COPY; X-spiders act like XOR:


Classically we have:


## Bialgebra

Z-spiders act like COPY; X-spiders act like XOR:


Classically we have:


Hence:


Three CNOTs make SWAP


Three CNOTs make SWAP


Three CNOTs make SWAP


Three CNOTs make SWAP


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## Rules for ZX-diagrams: The ZX-calculus



$$
\cdots-\alpha=\underset{\vdots}{0-} \quad \alpha, \beta \in[0,2 \pi]
$$

- All derivations hold in any orientation
- All derivations hold with colours interchanged
- All derivations hold with phases negated


## Example 1: GHZ-preparation circuit

Recall that the GHZ-state is $|000\rangle+|111\rangle$. The following circuit creates a GHZ-state:


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## Example 2: Teleportation

Let $|\Psi\rangle$ represent a side of a Bell state.
Then this is the standard quantum teleportation protocol:


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Now let's look at specific use-cases of ZX

## Use-case of ZX \#1: Clifford computation

Gottesman-Knill theorem
A quantum circuit of Cliffords can be efficiently classically simulated.

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Gottesman-Knill theorem
A quantum circuit of Cliffords can be efficiently classically simulated.
Can we prove this using ZX ?

## Cliffords in ZX

- A Clifford map is any linear map produced from combining Clifford unitaries, states and post-selections $\langle 0|$.


## Cliffords in ZX

- A Clifford map is any linear map produced from combining Clifford unitaries, states and post-selections $\langle 0|$.
- As a ZX-diagram, a Clifford map only has phases multiple of $\frac{\pi}{2}$.

$$
\mathrm{CNOT}=\underset{0}{0-\left(\frac{\pi}{2}\right.} \quad \mathrm{S}=-\left(\frac{\pi}{2}-\left(\frac{\pi}{2}\right)-\left(\frac{\pi}{2}\right)-\right.
$$

- Conversely, ZX-diagrams with phases multiple of $\frac{\pi}{2}$ are Clifford.


## Graph-like diagrams

Every ZX-diagram can be reduced to a graph-like diagram:

- Only Z-spiders and Hadamards:



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## Graph-like diagrams

Every ZX-diagram can be reduced to a graph-like diagram:

- Only Z-spiders and Hadamards:

- Cancel adjacent hadamards: -
- Fuse all spiders.
- No self-loops or multiple edges:

- View all Hadamards as a type of edge:



## Graph states

A graph-like diagram is a graph state when

- it has no inputs,
- every spider is connected to a unique output,
- all phases are zero.

Example:


## Graph-theoretic rewriting

We've transformed ZX-diagrams into simple undirected graphs so we can view rewrites graph-theoretically.

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Local complementation
G


$G \star a:=G$, but with connectivity of neighbours of a complemented.
A pivot on edge $u v$ is $G \wedge u v:=G \star u \star v \star u$.

G

$G \wedge u v$


## Local complementation on graph states

An Icomp on graph state can be implemented using local Cliffords:


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An Icomp on graph state can be implemented using local Cliffords:


Same goes for pivot:


## Removing vertices

Remove vertex by Icomp:


Similarly, using pivot:


## Clifford simplification

- With Icomp can remove all internal vertices with $\pm \frac{\pi}{2}$ phase.
- With pivot can remove all internal vertices with 0 or $\pi$ phase.


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Gottesman-Knill theorem
A Clifford computation can be efficiently classically simulated.

## Clifford normal form

Other consequence: Clifford circuit reduced to


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Normal form of layers:

$$
\mathrm{H}+\mathrm{S}+\mathrm{CZ}+\mathrm{CNOT}+\mathrm{H}+\mathrm{CZ}+\mathrm{S}+\mathrm{H}
$$

## Simplifying general circuits

Example result after simplification:


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Problem: does not look a circuit.

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Example result after simplification:


Problem: does not look a circuit.
Solution: all rewrites preserve gflow.

- Duncan, Perdrix, Kissinger, vdW (2019). Graph-theoretic Simplification of Quantum Circuits with the ZX-calculus.
- Backens, Miller-Bakewell, de Felice, Lobski, vdW (2020). There and back again: A circuit extraction tale.

Non-Clifford optimisation

## Non-Clifford optimisation

Additional rules for phase gadgets:


$$
:: \quad\left|x_{1}, \ldots, x_{n}\right\rangle \mapsto e^{i \alpha\left(x_{1} \oplus \ldots \oplus x_{n}\right)}\left|x_{1}, \ldots, x_{n}\right\rangle
$$

## Non-Clifford optimisation

Additional rules for phase gadgets:


$$
\therefore \quad\left|x_{1}, \ldots, x_{n}\right\rangle \mapsto e^{i \alpha\left(x_{1} \oplus \ldots \oplus x_{n}\right)}\left|x_{1}, \ldots, x_{n}\right\rangle
$$



Kissinger, vdW 2019: Reducing T-count with the ZX-calculus

## T-count optimisation

- Phase gadget optimisation allows us to kill non-Clifford phases.


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- Combining with TODD [Heyfron \& Campbell 2018] we improved T-counts for 20/36 circuits.
- Note: [Zhang \& Chen 2019] use a different method that achieves nearly identical T-counts.


## Circuit equality verification

Can we verify correctness of optimisations?

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- Compose optimised circuit with adjoint of original circuit
- Simplify
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- Compose optimised circuit with adjoint of original circuit
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Using this method found mistake in other peer-reviewed optimiser.

## Classical Quantum circuit optimisation

Two ways for ZX to classically simulate circuits:

- Treat ZX-diagram as tensor network and contract.


## Classical Quantum circuit optimisation

Two ways for ZX to classically simulate circuits:

- Treat ZX-diagram as tensor network and contract.
- Use stabiliser decomposition of magic states to write as sum of simpler diagrams.


FIG. 3. Graphs $G^{\prime}$ and $G^{\prime \prime}$ used in the definition of stabilizer states $\phi^{\prime}$ and $\phi^{\prime \prime}$; see Eq. (11).

$$
\begin{align*}
\left|H^{\otimes 6}\right\rangle= & (-16+12 \sqrt{2})\left|B_{6,0}\right\rangle+(96-68 \sqrt{2})\left|B_{6,6}\right\rangle \\
& +(10-7 \sqrt{2})\left|E_{6}\right\rangle+(-14+10 \sqrt{2})\left|O_{6}\right\rangle \\
& +(7-5 \sqrt{2}) Z^{\otimes 6}\left|K_{6}\right\rangle+(10-7 \sqrt{2})\left|\phi^{\prime}\right\rangle \\
& +(10-7 \sqrt{2})\left|\phi^{\prime \prime}\right\rangle, \tag{11}
\end{align*}
$$

where

$$
\left|\phi^{\prime}\right\rangle=\prod_{(i, j) \in E^{\prime}} \Lambda(Z)_{i, j}\left|O_{6}\right\rangle \text { and }\left|\phi^{\prime \prime}\right\rangle=\prod_{(i, j) \in E^{\prime \prime}} \Lambda(Z)_{i, j}\left|O_{6}\right\rangle .
$$

Source: Sergey Bravyi, Graeme Smith, and John A Smolin. Trading classical and quantum computational resources (2016).

$$
\begin{aligned}
& e^{i \pi / 4} \underset{(1)}{\frac{\pi}{4}}\left(\frac{\pi}{4}\right)\left(\frac{\pi}{4}\right)\left(\frac{\pi}{4}\left(\frac{\pi}{4}\right)\left(\frac{\pi}{4}\right)\right. \\
& -\frac{1+\sqrt{2}}{4} 0+2 e^{i \pi / 4} \\
& -2 \sqrt{2} i
\end{aligned}
$$

## Circuit simulation with ZX-calculus

1. Write circuit+state as $Z X$-diagram.
2. Simplify using $Z X$-calculus rules.
3. Replace magic states by stabilizer decomposition.
4. Repeat.
5. ...
6. Profit!

Early results looks like this could give major benefit

## Stuff I didn't talk about

- CNOT optimisation
- Relationship to MBQC and lattice surgery
- Circuit routing
- Applications in tensor networks


## Conclusion

- ZX-calculus is a better representation of quantum circuits
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## Thank you for your attention

vdW 2020, arXiv:2012.13966.
ZX-calculus for the working quantum computer scientist

