



Quantum theory is a quasi-stochastic process theory

Radboud University

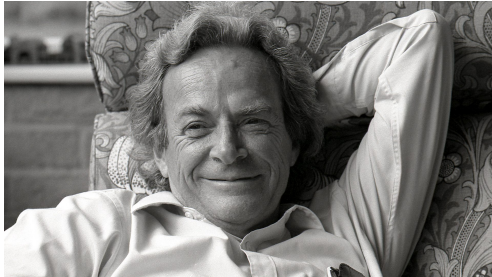
John van de Wetering

wetering@cs.ru.nl

Institute for Computing and Information Sciences
Radboud University Nijmegen

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“The only difference between a probabilistic classical world and the equations of the quantum world is that somehow or other it appears as if the probabilities would have to go negative.”

– Richard Feynman, 1981



A bit of background

- Wigner (1932): Representing a quantum state as a distribution over classical phase space allowing negative probabilities.





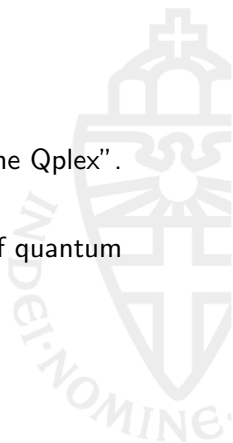
A bit of background

- Wigner (1932): Representing a quantum state as a distribution over classical phase space allowing negative probabilities.
- Negativity in representations is “equivalent” to contextuality (Spekkens 2008).
- Quantum speed up requires sufficient negativity in representations (Pashayan, Walman & Bartlett 2015).



Related work

- Appleby, Fuchs, Stacey, Zhu 2016 “Introducing the Qplex”.
- Hardy 2013 “The duotensor framework”
- Ferrie & Emerson 2008 “Frame representations of quantum mechanics”





Informationally complete POVMs

Definition

- Let M_n be the set of $n \times n$ complex matrices.
- An *effect* is an $E \in M_n$ such that $0 \leq E \leq 1$.
- A *POVM* is a set of effects $\{E_i\}$ such that $\sum_i E_i = I_n$.



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- A POVM is called *informationally complete* if it spans M_n and *minimal* informationally complete (MIC) if it is a basis. A MIC-POVM always has n^2 elements.



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Definition

A *quantum state* is $\rho \in M_n$ such that $\rho \geq 0$ and $\text{tr}(\rho) = 1$.



Quantum states as probability distributions

Let $\rho \in M_n$ be a quantum state and $\{E_i\}$ a POVM.
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Now:

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Quantum states as probability distributions - cont.

$$p(i) = \text{tr}(\rho E_i) = \sum_j \alpha_j \text{tr}\left(\frac{E_j}{\text{tr}(E_j)} E_i\right)$$

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then we can succinctly write

$$p = T\alpha \quad \text{or equivalently} \quad \alpha = T^{-1}p$$

Which allows us to reconstruct the original state:

$$\rho = \sum_i (T^{-1}p)_i \frac{E_i}{\text{tr}(E_i)}$$

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NOTE: T^{-1} can contain negative components!



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- A real-valued matrix S is called *stochastic* when $S_{ij} \in \mathbb{R}_{\geq 0}$ for all i, j and all the columns sum up to 1.
- It is *quasi-stochastic* when the positivity requirement is dropped.
- S is *doubly* (quasi-)stochastic when its transpose is also (quasi-)stochastic.

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S stochastic $\not\Rightarrow S^{-1}$ stochastic (when it exists).



Quantum channels as quasi-stochastic matrices

Let $\Phi : M_n \rightarrow M_m$ be a CPTP-map and fix MIC-POVMs $\{E_i\}$ and $\{E'_j\}$ on respectively M_n and M_m . Let T be the transition matrix for $\{E_i\}$.





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$$\rho(i) := \text{tr}(\rho E_i) \quad \Rightarrow \quad \rho = \sum_i (T^{-1} \rho)_i \frac{E_i}{\text{tr}(E_i)}$$

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Define

$$Q(\Phi)_{ij} = \text{tr}\left(\Phi\left(\frac{E_j}{\text{tr}(E_j)}\right) E'_i\right)$$

Then $q = Q(\Phi)T^{-1}p$



Now that we've got that out of the way...

...time for some new stuff





Quantum channels as quasi-stochastic matrices - cont.

$\Phi : M_n \rightarrow M_m$ and $\Psi : M_m \rightarrow M_k$ with MIC-POVMs $\{E_i\}$, $\{E'_i\}$ and $\{E''_i\}$, and transition matrices T , T' and T'' .





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and set $\tau = (\Psi \circ \Phi)(\rho)$ with distribution $r(i) = \text{tr}(\tau E''_i)$.



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Quantum theory as a quasi-stochastic process theory

Definition

Let **CPTP** be the category with objects natural numbers and morphisms CPTP maps $\Phi : M_n \rightarrow M_m$.

Let **QStoch** be the category with objects natural numbers and morphisms quasi-stochastic matrices.

Note: Density matrices are equivalent to $\hat{\rho} : M_1 = \mathbb{C} \rightarrow M_n$.



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Fix $\forall n \in \mathbb{N}$ MIC-POVMs $\{E_i^{(n)}\}$ with transition matrices T_n .

Let $F_E : \mathbf{CPTP} \rightarrow \mathbf{QStoch}$ be a functor with $F_E(n) = n^2$ and $F_E(\Phi : M_n \rightarrow M_m) = Q(\Phi)T_n^{-1}$ where

$$Q(\Phi)_{ij} = \text{tr} \left(\Phi \left(\frac{E_j^{(n)}}{\text{tr}(E_j^{(n)})} \right) E_i^{(m)} \right)$$



Properties of the quasi-stochastic representation

Theorem

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A different set of MIC-POVMs gives a different functor, but:

Theorem

Any two functors $F_E, F_{E'} : \mathbf{CPTP} \rightarrow \mathbf{QStoch}$ arising from a choice of MIC-POVMs are naturally isomorphic.



Preservation of tensor product

Definition: Strong monoidal functors

A functor $F : \mathbb{A} \rightarrow \mathbb{B}$ is called *strong monoidal* if there exist isomorphisms $\alpha_{A,B}$ for every pair of objects A and B such that $\alpha_{B_1, B_2} \circ (F(f_1) \otimes F(f_2)) = F(f_1 \otimes f_2) \circ \alpha_{A_1, A_2}$ for all morphisms $f_i : A_i \rightarrow B_i$ satisfying some coherence conditions.



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Theorem

The functor $F_E : \mathbf{CPTP} \rightarrow \mathbf{QStoch}$ is strong monoidal.

NOTE: You need minimality of the POVMs for this!



Preservation of adjoints

Definition: Linear algebraic adjoint

Let $A : (V, \langle \cdot, \cdot \rangle) \rightarrow (W, \langle \cdot, \cdot \rangle)$ be a linear map. Its *adjoint* is a map $A^\dagger : (W, \langle \cdot, \cdot \rangle) \rightarrow (V, \langle \cdot, \cdot \rangle)$ such that

$$\langle v, A^\dagger w \rangle = \langle Av, w \rangle$$

e.g. adjoint of real matrix is the transpose
and adjoint of $\hat{U}(A) = UAU^\dagger$ is $\hat{U}^\dagger(A) = U^\dagger AU$.

The adjoint of a CPTP map is CPTP if and only if it is *unital*.



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Answer: No! (in general)



Symmetric Informationally Complete POVMs

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A MIC-POVM $\{E_i\}$ is called *symmetric* when

$$\exists \alpha, \beta : \forall i, j : \text{tr}(E_i E_j) = \alpha \delta_{ij} + \beta$$

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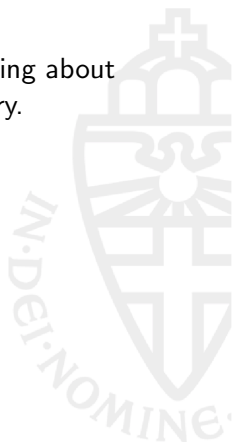
Theorem

The functor $F_E : \mathbf{CPTP} \rightarrow \mathbf{QStoch}$ preserves the adjoint of unital channels, e.g. $F(\Phi^\dagger) = F(\Phi)^\dagger$, if and only if all associated MIC-POVMS are symmetric.



Conclusion and Discussion

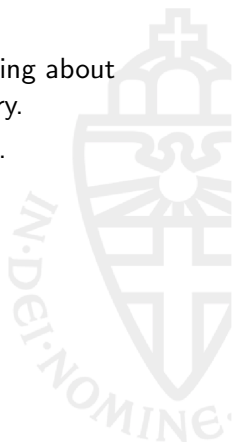
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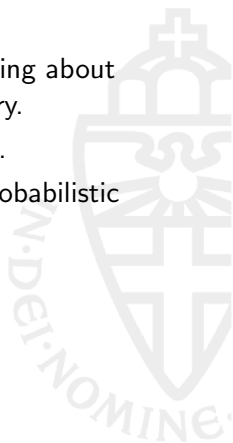
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- Yet again a special role for symmetric IC-POVMs.
- Construction also applies to causal operational probabilistic theories.
- **QStoch** doesn't 'care' about positivity. Can this be fixed?
- Can we 'simulate' causal OPTs using quantum theory with these representations?



Thank you for your attention

